

# 2D COMPUTER GRAPHICS

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Diego Nehab

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IMPA

# DIFFERENTIAL GEOMETRY

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## INSIDE-OUTSIDE TEST

Ideal model is *implicit*

Given a region  $\Omega \subset \mathbf{R}^2$ , define the *indicator function*  $\mathbf{1}_\Omega : \mathbf{R}^2 \rightarrow \{0, 1\}$

$$\mathbf{1}_\Omega(p) = \begin{cases} 1, & p \in \Omega \\ 0, & p \notin \Omega \end{cases}$$

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A.k.a. *characteristic function*  $\chi_\Omega(p) = \mathbf{1}_\Omega(p)$

Alternatively, using the *Iverson bracket*,  $[p \in \Omega] = \mathbf{1}_\Omega(p)$ , where  $[true] = 1$  and  $[false] = 0$

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Basis of CSG (constructive solid geometry)

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How do we define the boundary?

## PLANAR PARAMETRIC CURVE

Piecewise differentiable function  $\alpha : I \subset \mathbf{R} \rightarrow \mathbf{R}^2$  from an interval  $I$  to  $\mathbf{R}^2$

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A subset  $S \subset \mathbf{R}^2$  can be parametrized in many different ways

$$\beta : [a, b] \rightarrow \mathbf{R}^2 \quad \beta(t) = (r \cos(\omega t + \phi), r \sin(\omega t + \phi)), \quad b - a \geq \frac{2\pi}{\omega}$$

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Are individual SVG *segments* regular?

## ARC LENGTH

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- Makes sense from physics' time integral of speed
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- Koch snowflake



## REPARAMETERIZATION

A curve  $\beta : J \rightarrow \mathbf{R}^2$  is a *reparameterization* of  $\alpha$  if there is a monotonic differentiable function  $h : J \rightarrow I$  such that  $\beta = \alpha \circ h$

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$$\begin{aligned}\int_c^d |\beta'(t)| dt &= \int_c^d |\alpha'(h(t))| h'(t) dt \\ &= \int_a^b |\alpha'(u)| du \quad (u = h(t))\end{aligned}$$



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Regularity means  $s'(t) = |\alpha'(t)| > 0$ , which means  $s(t)$  is strictly increasing, which means  $s$  has a differentiable inverse  $u$  with

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We get

$$s(t) = \int_c^t |\beta'(t)| dt = t - c$$

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$$\alpha(t) = (at^2, 2at) \Rightarrow$$

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Standard ellipse

$$\beta(t) = (a \cos t, b \sin t) \Rightarrow$$

$$|\beta'(t)| = b\sqrt{1 - m \sin^2(t)}, \quad m = 1 - \frac{a^2}{b^2}$$

$$\int_0^t |\beta'(t)| dt = \text{Elliptic integral of the second kind}$$

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$N(t) = T'(t)/|T'(t)|$  is the *unit normal* to  $\beta$



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An inflection is a point where the curvature vanishes

I.e. where the 1st and 2nd derivatives are collinear

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## BÉZIER CURVES

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Show offset and evolute curves



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